



Answers to Practice Test 4

Proofs by Induction. Use your blue books.

$$8. \sum_{i=1}^n i(i!) = (n+1)! - 1.$$

Base Case. Let $n = 1$. It suffices to prove that $1 \cdot 1! = (1 + 1)! - 1$
 $1 \cdot 1! = 1 \cdot 1$
 $= 1$
 $= 2 - 1$
 $= 2! - 1$
 $= (1 + 1)! - 1$

Inductive Case.
 Inductive hypothesis: the sum from 1 to $k = (k + 1)! - 1$
 It suffices to show that the sum from 1 to $k + 1 = (k + 2)! - 1$.
 Sum from 1 to $k + 1 =$ sum from 1 to $k + (k+1)(k+1)!$
 $= (k + 1)! - 1 + (k+1)(k+1)!$, by IH.
 $= (k + 2) \cdot (k+1)! - 1$

9. Consider the formal language consisting of vocabulary items  and . Grammatical strings are defined recursively as follows:

- (1) $\text{List}(\text{book icon})$ is a string.
- (2) If \mathcal{A} is a string, then $\text{List}(\text{book icon})^{\mathcal{A}} \text{List}(\text{box icon})$ is a string.
- (3) There are no other strings.

Prove that in any string, the number of  's is greater than the number of  s.

First, we use recursion to define two functions, the number of k 's function (n) and the number of $=$'s function (q):

$$n(\text{List}(k)) = 1$$

$$n(\text{List}(k)^{\mathcal{A}} \text{List}(=)) = n(\mathcal{A}) + 1$$

$$q(\text{List}(k)) = 0$$

$$q(\text{List}(k)^{\mathcal{A}} \text{List}(=)) = q(\mathcal{A}) + 1$$

Let φ be an arbitrary string
 Base Case. String $\varphi = \text{List}(k)$.
 $n(\varphi) = 1, q(\varphi) = 0$, and $1 > 0$.

Inductive case. $\varphi = \text{List}(k)^{\psi} \text{List}(=)$.
 Inductive hypothesis: $n(\psi) > q(\psi)$.
 $n(\varphi) = n(\psi) + 1$, and $q(\varphi) = q(\psi) + 1$,
 So $n(\varphi) > q(\varphi)$.

Bonus Problem. (10 points maximum)

Use the following recursive definition of the concatenation function on lists :

1. If x is a list, then $x^{\wedge}NIL = x$.
2. If x and y are lists, then $x^{\wedge}Cons(y,a) = Cons(x^{\wedge}y, a)$.

Prove that for all lists x, y and z , $(x^{\wedge}y)^{\wedge}z = x^{\wedge}(y^{\wedge}z)$. [Hint: prove by induction on the length of z .]

Base Case. Let $z = NIL$. Show for all x, y , $(x^{\wedge}y)^{\wedge}NIL = x^{\wedge}(y^{\wedge}NIL)$.
 $(x^{\wedge}y)^{\wedge}NIL = (x^{\wedge}y)$.
 $= x^{\wedge}(y^{\wedge}NIL)$, since $y = y^{\wedge}NIL$.

Inductive Case. Let $z = Cons(w,a)$.
Inductive hypothesis: for all x, y , $(x^{\wedge}y)^{\wedge}w = x^{\wedge}(y^{\wedge}w)$.
Show: for all x, y , $(x^{\wedge}y)^{\wedge}z = x^{\wedge}(y^{\wedge}z)$,
 $(x^{\wedge}y)^{\wedge}z$
 $= (x^{\wedge}y)^{\wedge}Cons(w,a)$
 $= Cons(((x^{\wedge}y)^{\wedge}w),a)$, by definition of \wedge
 $= Cons((x^{\wedge}(y^{\wedge}w)),a)$, by IH.
 $= (x^{\wedge}Cons((y^{\wedge}w),a))$, by definition of \wedge .
 $= (x^{\wedge}(y^{\wedge}Cons(w,a)))$, by definition of \wedge .
 $= (x^{\wedge}(y^{\wedge}z))$