

Practice Exam 4

I. Multiple Choice.

1. Which of the following binary relations (on \mathbb{N}) is **not** a function from \mathbb{N} into \mathbb{N} ? [Think of x as the input and y as the output.]

- (a) $R_{xy} \leftrightarrow y = x + 3$
- (b) $S_{xy} \leftrightarrow y = x - 1$
- (c) $T_{xy} \leftrightarrow y = x$
- (d) $U_{xy} \leftrightarrow y = x^2 + 5$
- (e) $V_{xy} \leftrightarrow y = 2^x$

2. If $f: A \rightarrow B$ is bijective, then which is **not** necessarily true of f^{-1} ?

- (a) $f^{-1}: B \rightarrow A$ is bijective.
- (b) $\forall x \in A \quad f^{-1}(f(x)) = x$.
- (c) $\forall x \in B \quad f(f^{-1}(x)) = x$.
- (d) $f = f^{-1}$.
- (e) $f = (f^{-1})^{-1}$.

3. Which of the following functions is surjective (onto)?

- (a) $f: \mathbb{N} \rightarrow \mathbb{N}$, where $f(x) = x + 2$.
- (b) $g: \mathbb{N} \rightarrow \mathbb{N}$, where $g(x) = 3 \cdot x - 1$.
- (c) $h: \mathbb{Z} \rightarrow \mathbb{Z}$, where $h(x) = -x$.
- (d) $j: \mathbb{Z} \rightarrow \mathbb{Z}$, where $j(x) = x^2$.
- (e) $k: \mathbb{R} \rightarrow \mathbb{R}$, where $k(x) = x^2$.

4. Which of the following is **not** true of \mathbb{N} (the set of natural numbers)?

- (a) \mathbb{N} is inductive.
- (b) \mathbb{N} is infinite.
- (c) $<$ on \mathbb{N} is well-founded.
- (d) Every nonempty subset of \mathbb{N} has a least member.
- (e) Every member of \mathbb{N} is the successor of some member of \mathbb{N} .

5. What is wrong with the following would-be proof?

[Assume the following definition : y is odd $\leftrightarrow \exists z(z \in \mathbb{N} \ \& \ y = 2z + 1)$.]

"Show: for all natural numbers x , if $x \geq 1$, then $x^2 + x$ is odd.

By induction on x , starting with $x = 1$.

Base case: Show: 1 is odd.

$1 = 2 \cdot 0 + 1$, so 1 is odd, by definition.

Inductive hypothesis: Assume $n^2 + n$ is odd. That is, for some m , $n^2 + n = 2m + 1$

Show: $(n+1)^2 + (n+1)$ is odd.

$$\begin{aligned} (n+1)^2 + (n+1) &= (n^2 + 2n + 1) + (n + 1) \\ &= (n^2 + n) + (2n + 2) \\ &= (n^2 + n) + 2 \cdot (n + 1) \end{aligned}$$

By inductive hypothesis, $(n^2 + n) = 2 \cdot m + 1$, for some m . Therefore:

$$\begin{aligned} (n + 1)^2 + (n + 1) &= (2 \cdot m + 1) + 2 \cdot (n + 1) \\ &= 2 \cdot (m + n + 1) + 1 \end{aligned}$$

Therefore, $(n+1)^2 + (n+1)$ is odd, by definition."

- (a) Error in the statement of the base case.
- (b) Error in the proof of the base case.
- (b) Error in the statement of the inductive hypothesis.
- (c) Error in the statement of the inductive case.
- (d) Error in the proof of the inductive case.

6. What's wrong with the following would-be proof?

"Show: for all numbers x, y and z , $(x + y)z = xz + yz$.

By induction on x , starting with $x = 0$.

Base case: Show: for all numbers x and y , $(0 + y)z = 0 \cdot z + yz$.

By Cor. 38.1, $0 + y = y$, so $(0 + y) \cdot z = y \cdot z$. By Cor 38.1 again, $yz = 0 + yz$. By Cor. 39.1, $0z = 0$. So, $0 + y \cdot z = 0 \cdot z + y \cdot z$.

Inductive hypothesis: for all numbers y and z , $(n+y)z = nz + yz$.

Show: for all numbers y and z , $((n+1) + y)z = (n+1)y + yz$.

$$\begin{aligned} ((n+1) + y)z &= ((n+y) + 1)z, \text{ by Cor 38.2,} \\ &= (n+y)z + z, \text{ by Cor. 39.2.} \\ &= (nz + yz) + z, \text{ by IH.} \\ &= (nz + z) + yz, \text{ by commutativity and associativity of +.} \\ &= (n+1)z + yz, \text{ by Cor. 39.2.} \end{aligned}$$

Answers:

- (a) Error in Base case.
- (b) Error in statement of inductive hypothesis.
- (c) Error in statement of the Show line for inductive step.
- (d) Error in proof of inductive step.
- (e) No serious errors.

7. Consider the following definition of 2^{*n} , "2 to the super-power of n ":

$$2^{*0} = 2$$

$$2^{*(k+1)} = 2^{2^{*k}}$$

Which of the following correctly lists the first four values of this function (i.e.,

$2^{*0}, 2^{*1}, 2^{*2}, 2^{*3}$)?

- (a) 2, 4, 16, 256
- (b) 2, 4, 16, 128
- (c) 2, 4, 16, 64
- (d) 2, 4, 16, 65536
- (e) 0, 2, 4, 16

II. Proofs by Induction. Use your blue books.

$$8. \sum_{i=1}^n i(i!) = (n+1)! - 1.$$

9. Consider the formal language consisting of vocabulary items $\langle \text{List} \rangle$ and $\langle \text{Cons} \rangle$. Grammatical strings are defined recursively as follows:

- (1) $\langle \text{List} \rangle$ is a string.
- (2) If A is a string, then $\langle \text{List} \rangle _ A _ \langle \text{Cons} \rangle$ is a string.
- (3) There are no other strings.

Prove that in any string, the number of $\langle \text{List} \rangle$'s is greater than the number of $\langle \text{Cons} \rangle$'s.

Bonus Problem. (10 points maximum)

Use the following recursive definition of the concatenation function on lists :

1. If x is a list, then $x \wedge \text{NIL} = x$.
2. If x and y are lists, then $x \wedge \text{Cons}(y, a) = \text{Cons}(x \wedge y, a)$.

Prove that for all lists x, y and z , $(x \wedge y) \wedge z = x \wedge (y \wedge z)$. [Hint: prove by induction on the length of z .]